

# NOTES ON HIGHER-DIMENSIONAL TARAI FUNCTIONS

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## 0. INTRODUCTION

I. Takeuchi defined the following recursive function, called the tarai function, in [5].

$$t(x, y, z) = \text{if } x \leq y \text{ then } y \text{ else } t(t(x-1, y, z), t(y-1, z, x), t(z-1, x, y))$$

This function requires many recursive calls even for small  $x$ ,  $y$ , and  $z$ , so it is used to see how effectively the programming language implementation handles recursive calls. In [3], J. McCarthy proved that this recursion terminates without call-by-need and  $t$  can be computed in the following way.

$$t(x, y, z) = \text{if } x \leq y \text{ then } y \text{ else if } y \leq z \text{ then } z \text{ else } x$$

In [4], J. S. Moore gave an easier proof and verified it by the Boyer-Moore theorem prover.

D. Knuth proposed the following generalization in [2], called the  $n$ -dimensional tarai function.

$$t(x_1, x_2, \dots, x_n) = \text{if } x_1 \leq x_2 \text{ then } x_2 \\ \text{else } t(t(x_1-1, x_2, \dots, x_n), \dots, t(x_n-1, x_1, \dots, x_{n-1}))$$

It was shown by T. Bailey, J. Coldwell, and J. Cowles that the 4-dimensional tarai function does not terminate without call-by-need because

$$\begin{aligned} t(3, 2, 1, 5) &= t(t(2, 2, 1, 5), t(1, 1, 5, 3), t(0, 5, 3, 2), t(4, 3, 2, 1)) \\ &= t(2, 1, 5, 4) \\ &= t(t(1, 1, 5, 4), t(0, 5, 4, 2), t(4, 4, 2, 1), t(3, 2, 1, 5)) \end{aligned}$$

T. Bailey and J. Cowles announced in [1] that they gave an informal (handwritten) proof of the following conjecture.

**Conjecture 1.** Let  $n \geq 3$  be an integer. Define the function  $f$  on  $\mathbb{Z}^n$  by

$$\begin{aligned} f(x_1, x_2, \dots, x_m) &= \text{if } (\exists k < m)(x_1 > x_2 > \dots > x_k \leq x_{k+1}) \\ &\quad \text{then } g_b(x_1, x_2, \dots, x_{k+1}) \\ &\quad \text{else } x_1 \end{aligned}$$

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where the function  $g_b$  is defined by

$$g_b(x_1, x_2, \dots, x_j) = \begin{array}{ll} \text{if } j \leq 3 \text{ then } x_j & \\ \text{else if } x_1 = x_2 + 1 \text{ or } x_2 > x_3 + 1 & \\ \quad \text{then } g_b(x_2, \dots, x_j) & \\ \text{else } \max\{x_3, x_j\} & \end{array}$$

Then,  $f$  satisfies the  $n$ -dimensional tarai recurrence.

The goal of this paper is to give a proof to this theorem. Moreover, the proof will be simpler than the one proposed in [1], so we hope that it is easier to be formalized.

### 1. TERMINATION WITH CALL-BY-NEED

In this section, we shall prove that the  $n$ -dimensional tarai function is a total function for every  $n \geq 3$ . Throughout this section, let  $n$  be a fixed natural number with  $n \geq 3$  and  $t$  the  $n$ -dimensional tarai function.

First, we shall prepare notation. Let  $\vec{x} = \langle x_1, \dots, x_n \rangle \in \mathbb{Z}^n$ . Define  $\sigma(\vec{x})$  and  $r(\vec{x})$  by

$$\begin{aligned} \sigma(\vec{x}) &= \langle x_1 - 1, x_2, \dots, x_n \rangle \\ r(\vec{x}) &= \langle x_2, x_3, \dots, x_n, x_1 \rangle \end{aligned}$$

Namely, for every  $i \in \{1, \dots, n\}$ ,

$$\sigma(\vec{x})(i) = \begin{cases} \vec{x}(1) - 1 & \text{if } i = 1 \\ \vec{x}(i) & \text{otherwise} \end{cases}$$

and

$$r(\vec{x})(i) = \begin{cases} \vec{x}(i+1) & \text{if } i < n \\ \vec{x}(1) & \text{if } i = n \end{cases}$$

In particular, for every  $i \in \{1, \dots, n-1\}$  and  $j \in \{1, \dots, n\}$

$$r^i(\vec{x})(j) = \begin{cases} \vec{x}(j+i) & \text{if } j+i \leq n \\ \vec{x}(j+i-n) & \text{if } j+i > n \end{cases}$$

By using this notation, the  $n$ -dimensional tarai function  $t$  can be defined as for every  $\vec{x} \in \mathbb{Z}^n$ , if  $\vec{x}(1) \leq \vec{x}(2)$ , then  $t(\vec{x}) = \vec{x}(2)$ , and if  $\vec{x}(1) > \vec{x}(2)$ , then  $t(\vec{x}) = t(\vec{y})$ , where  $\vec{y} \in \mathbb{Z}^n$  is defined by  $\vec{y}(i) = t(\sigma(r^{i-1}(\vec{x})))$  for all  $i = 1, 2, \dots, n$ .

Let  $k \in \mathbb{N}$  with  $2 \leq k \leq n$ . Define  $X_k$  to be the set of all  $\vec{x} \in \mathbb{Z}^n$  such that for every  $i < k$ ,  $\vec{x}(i) \leq \vec{x}(k)$ . Note that the family  $\{X_k : 2 \leq k \leq n\}$  is not pairwise disjoint. For example,  $\langle 2, 1, 4, 3, 5 \rangle \in X_3 \cap X_5$ . Note also that if  $\vec{x} \in \mathbb{Z}^n$  satisfies  $\vec{x}(1) < \max \vec{x}$ , then there exists a  $k$  such that  $\vec{x} \in X_k$ .

**Lemma 1.1.** *For every  $k \in \mathbb{N}$  with  $2 \leq k \leq n$ , the following statement holds.*

- (\*)<sub>k</sub>: For every  $\vec{x} \in X_k$ ,
- (1)  $t(\vec{x})$  terminates with call-by-need,
  - (2)  $t(\vec{x})$  depends only on  $\vec{x} \upharpoonright \{1, \dots, k\}$ , and
  - (3)  $t(\vec{x}) \leq \vec{x}(k)$ .

*Proof.* Go by induction on  $k$ .

First assume  $k = 2$ . Then, we have  $\vec{x}(1) \leq \vec{x}(2)$  and hence  $t(\vec{x}) = \vec{x}(2)$ . Hence,  $\vec{x}$  clearly satisfies (i)–(iii).

Suppose that (\*)<sub>k'</sub> holds for all  $k' \in \mathbb{N}$  with  $2 \leq k' < k$ . We shall prove (\*)<sub>k</sub>. By way of contradiction, suppose that there exists an  $\vec{x} \in X_k$  which does not satisfy one of (i)–(iii). By inductive hypothesis, we have  $\vec{x} \notin X_{k-1}$ . In particular,  $\vec{x}(1) > \vec{x}(2)$ . Hence, we can pick the least  $y_1 \in \mathbb{Z}$  such that if  $\vec{y}$  is defined as  $\vec{y}(1) = y_1$  and  $\vec{y}(i) = \vec{x}(i)$  for every  $i \in \{2, \dots, n\}$ , then  $\vec{y}$  does not satisfy one of (i)–(iii). By redefining  $\vec{x}$ , we may assume that for every  $\vec{x}' \in \mathbb{N}^n$ , if  $\vec{x}'(1) < \vec{x}(1)$  and  $\vec{x}'(i) = \vec{x}(i)$  for every  $i \in \{2, \dots, n\}$ , then  $\vec{x}'$  satisfies (i)–(iii).

For each  $i \in \{1, \dots, k\}$ , define  $\vec{z}_i = \sigma(r^{i-1}(\vec{y}))$  and  $\vec{y}(i) = t(\vec{z}_i)$ . For  $i \in \{k+1, \dots, n\}$ , let  $\vec{y}(i)$  be left undefined. By definition,  $t(\vec{x})$  terminates with call-by-need if  $t(\vec{z}_i)$  for all  $i \in \{1, \dots, k\}$  and  $t(\vec{y})$  terminate with call-by-need, and in that case,  $t(\vec{x}) = t(\vec{y})$ .

Notice that for every  $i \in \{1, \dots, k-1\}$  and  $j \in \{2, \dots, k-i+1\}$ ,

$$\begin{aligned}\vec{z}_i(j) &= \vec{x}(j+i-1) \leq \vec{x}(k) \\ \vec{z}_i(1) &= \vec{x}(i) - 1 \leq \vec{x}(k)\end{aligned}$$

and

$$\vec{z}_i(k-i+1) = \vec{x}(i+k-i+1-1) = \vec{x}(k)$$

Hence,  $\vec{z}_i(k-i+1) = \vec{x}(k)$ . Therefore,  $\vec{z}_i \in X_{k-i+1}$ . If  $i \geq 2$ , then by inductive hypothesis, (\*)<sub>k-i+1</sub> holds. Thus,  $\vec{z}_i$  satisfies (i)–(iii). In particular,  $t(\vec{z}_i) \leq \vec{z}_i(k-i+1) = \vec{x}(k)$ . If  $i = 1$ , we have  $\vec{z}_1 = \sigma(\vec{x})$  and by the minimality of  $\vec{x}(1)$ , we know that  $\vec{z}_1$  satisfies (i)–(iii). For every  $i \in \{1, \dots, k-1\}$ ,  $\vec{y}(i) = t(\vec{z}_i) \leq \vec{x}(k)$ . Moreover,  $\vec{z}_{k-1}(2) = \vec{x}(k)$ . Thus,  $\vec{z}_{k-1}(1) \leq \vec{x}(k) = \vec{z}_{k-1}(2)$ , which implies  $\vec{y}(k-1) = t(\vec{z}_{k-1}) = \vec{z}_{k-1}(2) = \vec{x}(k)$ . Hence,  $\vec{y} \in X_{k-1}$  and so  $\vec{y}$  satisfies (i)–(iii). It is now easy to see that  $\vec{x}$  also satisfies (i)–(iii).  $\square$ (Lemma 1.1)

**Theorem 1.2.** For every  $\vec{x} \in \mathbb{Z}^n$ ,  $t(\vec{x})$  terminates with call-by-need and  $t(\vec{x}) \leq \max(\vec{x})$ .

*Proof.* By Lemma 1.1, it suffices to show that for every  $\vec{x} \in \mathbb{Z}^n$  with  $\vec{x}(1) = \max(\vec{x})$ ,  $t(\vec{x})$  terminates with call-by-need.

By way of contradiction, suppose that  $t(\vec{x})$  does not terminate. We may also assume that for every  $\vec{x}' \in \mathbb{Z}^n$  with  $\vec{x}'(1) < \vec{x}(1)$  and  $\vec{x}'(i) = \vec{x}(i)$  for every  $i \in \{2, \dots, n\}$ ,  $t(\vec{x}')$  terminates with call-by-need.

For every  $i \in \{1, \dots, n\}$ , let  $\vec{z}_i = \sigma(r^{i-1}(\vec{x}))$  and  $\vec{y}(i) = t(\vec{z}_i)$ . First suppose  $i \geq 2$ . Then, for every  $j \in \{2, \dots, n-i+1\}$ ,

$$\vec{z}_i(j) = \vec{x}(j+i-1) \leq \vec{x}(1)$$

and

$$\vec{z}_i(1) = \vec{x}(i) - 1 \leq \vec{x}(1)$$

Moreover,  $\vec{z}_i(n-i+2) = \vec{x}(n-i+2+(i-1)-n) = \vec{x}(1)$ . Thus,  $\vec{z}_i \in X_{n-i+2}$ . Hence,  $t(\vec{z}_i)$  terminates with call-by-need and  $t(\vec{z}_i) \leq \vec{x}(1)$ . In addition,  $\vec{z}_n(1) = \vec{x}(1+(n-1)) = \vec{x}(n) \leq \vec{x}(1)$  and  $\vec{z}_n(2) = \vec{x}(2+(n-1)-n) = \vec{x}(1)$ . Thus,  $t(\vec{z}_n) = \vec{x}(1)$ .

Suppose that  $i = 1$ . Then,  $\vec{z}_1 = \sigma(\vec{x})$ . By the minimality of  $\vec{x}(1)$ ,  $t(\vec{z}_1)$  terminates with call-by-need and  $t(\vec{z}_1) \leq \vec{x}(1)$ .

Therefore, for every  $i \in \{1, \dots, n\}$ , we have  $\vec{y}(i) = t(\vec{z}_i) \leq \vec{x}(1)$  and  $t(\vec{z}_n) = \vec{x}(1)$ . Hence, we have  $\vec{y} \in X_n$ . By  $(*)_n$ ,  $t(\vec{y})$  terminates with call-by-need and  $t(\vec{y}) \leq \vec{y}(n) = \vec{x}(1)$ . It follows that  $t(\vec{x})$  also terminates with call-by-need and  $t(\vec{x}) \leq \vec{x}(1)$ . This contradicts the choice of  $\vec{x}$ .  $\square$ (Theorem 1.2)

## 2. ALTERNATIVE DEFINITION OF THE FUNCTION OF T. BAILEY AND J. COWLES

In this section, we shall set up some definitions and notation which help us prove the main theorem. Let  $F$  denote the set of all non-empty finite sequences of integers.  $f$  denotes the function defined by T. Bailey and J. Cowles.

Let  $k$  be a function with domain  $F$  as follows. Let  $\vec{x} \in F$  be of length  $n$ . If  $\vec{x}(1) > \vec{x}(2) > \dots > \vec{x}(n)$ , then let  $k(\vec{x}) = n$ . Otherwise, let  $k(\vec{x})$  be the least  $k$  such that  $\vec{x}(k) \leq \vec{x}(k+1)$ .

Let  $l$  be a function with domain  $F$  as follows. Let  $\vec{x} \in F$  be of length  $n$ . If there is an integer  $l$  with  $1 \leq l < k(\vec{x})$  such that  $\vec{x}(l) > \vec{x}(l+1) + 1$  and  $\vec{x}(l+1) = \vec{x}(l+2) + 1$ , then let  $l(\vec{x})$  be the least such  $l$ . Otherwise, let  $l(\vec{x}) = k(\vec{x}) - 1$ . Notice that  $l(\vec{x}) = 0$  if and only if  $k(\vec{x}) = 1$ .

**Lemma 2.1.** *Let  $\vec{x} \in F$  be of length  $n \geq 3$ . Suppose that  $k(\vec{x}) = n - 1$ . Then,  $g_b(\vec{x}) = \max(\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1))$ .*

*Proof.* We shall prove the lemma by induction on  $n$ . If  $n = 3$ , then  $g_b(\vec{x}) = \vec{x}(3)$ . We also have  $k(\vec{x}) = 2$  and  $l(\vec{x}) = 1$ . Thus,  $\max(\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1)) = \vec{x}(3)$ . Therefore,  $g_b(\vec{x}) = \max(\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1))$ .

Suppose that the conclusion holds for all  $\vec{x}$  of length  $n$  for some  $n \geq 3$ . Let  $\vec{x} \in F$  be of length  $n+1$  with  $k(\vec{x}) = n$ . First suppose that  $\vec{x}(1) > \vec{x}(2) + 1$  and  $\vec{x}(2) = \vec{x}(3) + 1$ . Then  $g_b(\vec{x}) = \max(\vec{x}(3), \vec{x}(n+1))$ . It is clear that  $l(\vec{x}) = 1$ . Thus,  $\max(\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1)) = \max(\vec{x}(3), \vec{x}(n+1)) = g_b(\vec{x})$ .

Suppose  $\vec{x}(1) = \vec{x}(2) + 1$  or  $\vec{x}(2) > \vec{x}(3) + 1$ . Then  $g_b(\vec{x}) = g_b(\vec{y})$  where  $\vec{y}$  is a sequence of length  $n$  such that  $\vec{y}(i) = \vec{x}(i+1)$  for every  $i = 1, \dots, n$ .

Notice that  $k(\vec{y}) = k(\vec{x}) - 1$  and  $l(\vec{y}) = l(\vec{x}) - 1$ . By inductive hypothesis,  $g_b(\vec{y}) = \max(\vec{y}(l(\vec{y}) + 2), \vec{y}(k(\vec{y}) + 1)) = \max(\vec{y}(l(\vec{x}) + 1), \vec{y}(k(\vec{x}))) = \max(\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1))$ . Therefore,  $g_b(\vec{x}) = \max(\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1))$   $\square$ (Lemma 2.1)

By using the previous lemma, the following is immediate.

**Lemma 2.2.** *Let  $\vec{x} \in F$  be of length  $n \geq 3$ . If  $k(\vec{x}) = n$  (i.e.  $\vec{x}(1) > \vec{x}(2) > \dots > \vec{x}(n)$ ), then  $f(\vec{x}) = \vec{x}(1)$ . If  $k(\vec{x}) < n$ , then  $f(\vec{x}) = \max(\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1))$ .*

We shall use this characterization in the next section.

### 3. CLOSED FORM

In this section, we shall give a proof of the theorem of T. Bailey and J. Cowles. Note that if carefully rewritten, the proof simultaneously gives the termination of the  $n$ -dimensional tarai function. We chose to give separate proofs to simplify the arguments. Throughout this section, we fix a natural number  $n$  with  $n \geq 3$ .

We shall show that for every  $\vec{x} \in \mathbb{Z}^n$ ,  $t(\vec{x}) = f(\vec{x})$ . To this end, by Theorem 1.2, it suffices to show that  $f$  satisfies the  $n$ -dimensional tarai recurrence.

We shall begin with some easy facts about  $f$ .

**Lemma 3.1.** *For every  $\vec{x} \in \mathbb{Z}^n$ , if  $k(\vec{x}) \leq 2$ , then  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ .*

*Proof.* Since  $k(\vec{x}) \leq 2$ , we have  $l(\vec{x}) = k(\vec{x}) - 1$ . Thus,  $f(\vec{x}) = \max\{\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1)\} = \vec{x}(k(\vec{x}) + 1)$ .  $\square$ (Lemma 3.1)

**Lemma 3.2.** *For every  $\vec{x} \in \mathbb{Z}^n$  and  $m \in \{1, \dots, l(\vec{x}) + 2\}$ , if  $k(\vec{x}) < n$ , and  $\vec{x}(m) \leq \vec{x}(k(\vec{x}) + 1)$ , then  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ .*

*Proof.* If  $l(\vec{x}) = k(\vec{x}) - 1$ , then clearly  $f(\vec{x}) = \max\{\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1)\} = \vec{x}(k(\vec{x}) + 1)$ . Suppose  $l(\vec{x}) < k(\vec{x}) - 1$ . Then  $l(\vec{x}) + 2 \leq k(\vec{x})$ . We have  $m \leq l(\vec{x}) + 2 \leq k(\vec{x})$ . Thus,

$$\vec{x}(l(\vec{x}) + 2) \leq \vec{x}(m) \leq \vec{x}(k(\vec{x}) + 1)$$

So,  $f(\vec{x}) = \max\{\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1)\} = \vec{x}(k(\vec{x}) + 1)$ .  $\square$ (Lemma 3.2)

**Lemma 3.3.** *For every  $\vec{x} \in \mathbb{Z}^n$ , if  $k(\vec{x}) < n$  and  $\vec{x}(3) \leq \vec{x}(k(\vec{x}) + 1)$ , then  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ .*

*Proof.* By Lemma 3.1, if  $k(\vec{x}) \leq 2$ , then  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ . Suppose  $k(\vec{x}) \geq 3$ . Then by applying Lemma 3.2 with  $m = 3$ , we have  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ .  $\square$ (Lemma 3.3)

**Lemma 3.4.** *For every  $\vec{x} \in \mathbb{Z}^n$ , if  $k(\vec{x}) < n$  and  $\vec{x}(2) \leq \vec{x}(k(\vec{x}) + 1)$ , then  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ .*

*Proof.* If  $k(\vec{x}) \geq 3$ , then we have  $\vec{x}(3) < \vec{x}(2) \leq \vec{x}(k(\vec{x}) + 1)$ . By Lemma 3.3,  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ . If  $k(\vec{x}) \leq 2$ , then by Lemma 3.1, we have  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ .  $\square$ (Lemma 3.4)

**Lemma 3.5.** *For every  $\vec{x} \in \mathbb{Z}^n$ , if  $k(\vec{x}) < n$  and  $l(\vec{x}) \geq k(\vec{x}) - 2$ , then  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ .*

*Proof.* By definition,  $l(\vec{x}) \leq k(\vec{x}) - 1$ . So,  $l(\vec{x}) \geq k(\vec{x}) - 2$  implies either  $l(\vec{x}) = k(\vec{x}) - 2$  or  $l(\vec{x}) = k(\vec{x}) - 1$ . If  $l(\vec{x}) = k(\vec{x}) - 1$ , then clearly  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ . Suppose  $l(\vec{x}) = k(\vec{x}) - 2$ . Then,

$$\begin{aligned} f(\vec{x}) &= \max\{\vec{x}(l(\vec{x}) + 2), \vec{x}(k(\vec{x}) + 1)\} \\ &= \max\{\vec{x}(k(\vec{x})), \vec{x}(k(\vec{x}) + 1)\} \\ &= \vec{x}(k(\vec{x}) + 1) \end{aligned}$$

since by the definition of  $k(\vec{x})$ ,  $\vec{x}(k(\vec{x})) \leq \vec{x}(k(\vec{x}) + 1)$ .  $\square$ (Lemma 3.5)

**Lemma 3.6.** *For every  $\vec{x} \in \mathbb{Z}^n$ , if  $\vec{x}(2) \leq \vec{x}(3)$ , then  $f(\vec{x})$  is either  $\vec{x}(2)$  or  $\vec{x}(3)$ . In particular, if  $\vec{x}(2) = \vec{x}(3)$ , then  $f(\vec{x}) = \vec{x}(2)$ .*

*Proof.* Since  $\vec{x}(2) \leq \vec{x}(3)$ , we have  $k(\vec{x}) \leq 2$ . Then,  $k(\vec{x}) - 2 \leq 0 \leq l(\vec{x})$ . By Lemma 3.5,  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ . Since  $k(\vec{x})$  is either 1 or 2,  $f(\vec{x})$  is either  $\vec{x}(2)$  or  $\vec{x}(3)$ .  $\square$ (Lemma 3.6)

**Lemma 3.7.** *For every  $\vec{x} \in \mathbb{Z}^n$  and  $m \in \{1, \dots, n-1\}$ , if  $\vec{x}(m) = \vec{x}(m+1)$ ,  $k(\vec{x}) \geq m-1$ , and  $l(\vec{x}) \geq m-2$ , then  $f(\vec{x}) = \vec{x}(m)$ .*

*Proof.* Since  $\vec{x}(m) = \vec{x}(m+1)$ , we have  $k(\vec{x}) \leq m$ . So,

$$l(\vec{x}) \geq m-2 \geq k(\vec{x}) - 2$$

Hence, by Lemma 3.5,  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1)$ . However, since  $m-1 \leq k(\vec{x}) \leq m$ ,  $k(\vec{x})$  is either  $m-1$  or  $m$ . If  $k(\vec{x}) = m-1$ , then  $\vec{x}(k(\vec{x}) + 1) = \vec{x}(m)$ . If  $k(\vec{x}) = m$ , then  $\vec{x}(k(\vec{x}) + 1) = \vec{x}(m+1) = \vec{x}(m)$  by assumption.  $\square$ (Lemma 3.7)

**Lemma 3.8.** *For every  $\vec{x} \in \mathbb{Z}^n$  and  $m \in \{1, \dots, n-2\}$ , if  $k(\vec{x}) \geq m-1$ ,  $l(\vec{x}) \geq m-2$ ,  $\vec{x}(m) = \vec{x}(m+2)$ , and  $\vec{x}(m) \geq \vec{x}(m+1)$ , then  $f(\vec{x}) = \vec{x}(m)$ .*

*Proof.* By assumption,  $\vec{x}(m+1) \leq \vec{x}(m) = \vec{x}(m+2)$ . So,  $k(\vec{x}) \leq m+1$ . Therefore,  $k(\vec{x})$  is either  $m-1$ ,  $m$ , or  $m+1$ .

**Case 1.**  $k(\vec{x}) = m-1$ .

Then,  $l(\vec{x}) \geq m-2 = k(\vec{x}) - 1$ . By Lemma 3.5,  $f(\vec{x}) = \vec{x}(k(\vec{x}) + 1) = \vec{x}(m)$ .

**Case 2.**  $k(\vec{x}) = m$ .

In this case, we have  $\vec{x}(m) \leq \vec{x}(m+1)$ . By assumption, we also have  $\vec{x}(m+1) \leq \vec{x}(m)$ . Therefore,  $\vec{x}(m) = \vec{x}(m+1)$ . By Lemma 3.7,  $f(\vec{x}) = \vec{x}(m)$ .

**Case 3.**  $k(\vec{x}) = m+1$  and  $l(\vec{x}) = m-2$ .

Then,  $\vec{x}(k(\vec{x})+1) = \vec{x}(m+2) = \vec{x}(m)$  and  $\vec{x}(l(\vec{x})+2) = \vec{x}(m)$ . Thus,  $f(\vec{x}) = \vec{x}(m)$ .

**Case 4.**  $k(\vec{x}) = m+1$  and  $l(\vec{x}) \geq m-1$ .

Note that  $l(\vec{x}) \geq m-1 = k(\vec{x})-2$ . By Lemma 3.5,  $f(\vec{x}) = \vec{x}(k(\vec{x})+1) = \vec{x}(m+2) = \vec{x}(m)$ .  $\square$ (Lemma 3.8)

**Lemma 3.9.** *For every  $\vec{x} \in \mathbb{Z}^n$ , if  $2 \leq k(\vec{x}) < n$ , then  $f(\vec{x}) \leq \max\{\vec{x}(3), \vec{x}(k(\vec{x})+1)\}$ .*

*Proof.* If  $f(\vec{x}) = \vec{x}(k(\vec{x})+1)$ , it is trivial. Thus, we assume that  $f(\vec{x}) = \vec{x}(l(\vec{x})+2) \neq \vec{x}(k(\vec{x})+1)$ . It is easy to see that  $l(\vec{x}) \leq k(\vec{x})-2$ . So,  $l(\vec{x})+2 \leq k(\vec{x})$ . If  $l(\vec{x}) = 0$ , then  $k(\vec{x}) = 1$ , which contradicts  $l(\vec{x}) \leq k(\vec{x})-2$ . So, we have  $l(\vec{x}) \geq 1$ . Therefore,  $3 \leq l(\vec{x})+2 \leq k(\vec{x})$ . Hence,  $\vec{x}(3) \geq \vec{x}(l(\vec{x})+2) = f(\vec{x})$ .  $\square$ (Lemma 3.9)

**Lemma 3.10.** *For every  $\vec{x} \in \mathbb{Z}^n$ , if  $k(\vec{x}) < n$ , then  $f(\vec{x})$  satisfies the tarai recurrence.*

*Proof.* Let  $k = k(\vec{x})$  and  $l = l(\vec{x})$ . If  $k = 1$ , it is trivial. Assume  $k \geq 2$ .

For each  $i = 1, \dots, n$ , define  $\vec{z}_i = \sigma(r^{i-1}(\vec{x}))$  and let  $\vec{y} \in \mathbb{Z}^n$  be defined by  $\vec{y}(i) = f(\vec{z}_i)$ . It suffices to show  $f(\vec{x}) = f(\vec{y})$ .

It is easy to see that

$$\vec{z}_i(j) = \begin{cases} \vec{x}(i) - 1 & \text{if } j = 1 \\ \vec{x}(j+i-1) & \text{if } j+i-1 \leq n \\ \vec{x}(j+i-1-n) & \text{if } j+i-1 > n \end{cases}$$

Define  $m$  to be the least such that  $\vec{x}(m) > \vec{x}(m+1) + 1$  or  $m = k-1$ . Clearly we have  $m \leq l$ .

**Claim 1.**  $\vec{y}(k) = \vec{x}(k+1)$

$\vdash$  Note that  $\vec{z}_k(1) = \vec{x}(k)-1$  and  $\vec{z}_k(2) = \vec{x}(k+1)$ . Since  $\vec{x}(k) \leq \vec{x}(k+1)$ , we have  $\vec{z}_k(1) \leq \vec{z}_k(2)$ . Hence,  $\vec{y}(k) = f(\vec{z}_k) = \vec{z}_k(2) = \vec{x}(k+1)$ .  $\dashv$  (Claim 1)

**Claim 2.**  $k(\vec{y}) \leq k-1$ .

$\vdash$  Note  $\vec{z}_{k-1}(2) = \vec{x}(k)$  and  $\vec{z}_{k-1}(3) = \vec{x}(k+1)$ . So,  $\vec{z}_{k-1}(2) \leq \vec{z}_{k-1}(3)$ . By Lemma 3.6,  $\vec{y}(k-1)$  is either  $\vec{z}_{k-1}(2) = \vec{x}(k)$  or  $\vec{z}_{k-1}(3) = \vec{x}(k+1)$ . In either way, we have  $\vec{y}(k-1) \leq \vec{x}(k+1) = \vec{y}(k)$ . Thus,  $k(\vec{y}) \leq k-1$ .  $\dashv$  (Claim 2)

**Claim 3.** For every  $i \in \{1, \dots, n-1\}$ , if  $\vec{x}(i) = \vec{x}(i+1) + 1$ , then  $\vec{y}(i) = \vec{x}(i+1)$ . In particular, if  $i < m$ , then  $\vec{y}(i) = \vec{x}(i+1)$ .

⊢ We have  $\vec{z}_i(1) = \vec{x}(i) - 1$  and  $\vec{z}_i(2) = \vec{x}(i+1)$ . Thus,  $\vec{z}_i(1) = \vec{z}_i(2)$ .  
So,  $\vec{y}(i) = \vec{z}_i(2) = \vec{x}(i+1)$ . ⊣ (Claim 3)

**Claim 4.** For every  $i \in \{1, \dots, m-2\}$ ,  $\vec{y}(i) = \vec{y}(i+1) + 1$ . In particular,  $k(\vec{y}) \geq m-1$  and  $l(\vec{y}) \geq m-2$ . If  $l(\vec{y}) < k(\vec{y}) - 1$ , then  $l(\vec{y}) \geq m-1$ .

⊢ If  $i \in \{1, \dots, m-2\}$ , then by Claim 3,  $\vec{y}(i) = \vec{x}(i+1)$  and  $\vec{y}(i+1) = \vec{x}(i+2)$ . Since  $i+1 < m$ , we have  $\vec{x}(i+1) = \vec{x}(i+2) + 1$ . Thus,  $\vec{y}(i) = \vec{y}(i+1) + 1$ . So we have  $k(\vec{y}) \geq m-1$ . If  $l(\vec{y}) = k(\vec{y}) - 1$ , then  $l(\vec{y}) \geq m-1-1 = m-2$ . If  $l(\vec{y}) < k(\vec{y}) - 1$ , then we have  $\vec{y}(l(\vec{y})) > \vec{y}(l(\vec{y})+1) + 1$ . But we also have  $\vec{y}(m-2) = \vec{y}(m-1)$  and hence  $l(\vec{y}) \neq m-2$ . So,  $l(\vec{y}) \geq m-1$ . ⊣ (Claim 4)

**Claim 5.** For every  $i \in \{1, \dots, k-1\}$ , if  $\vec{x}(i) > \vec{x}(i+1) + 1$ , then  $k(\vec{z}_i) = k-i+1$  and  $\vec{z}_i(k(\vec{z}_i)+1) = \vec{x}(k+1)$ .

⊢ Note  $\vec{z}_i(1) = \vec{x}(i) - 1$  and  $\vec{z}_i(2) = \vec{x}(i+1)$ . By assumption,  $\vec{z}_i(1) > \vec{z}_i(2)$ . For every  $j \in \{2, \dots, k-i\}$ , we have  $\vec{z}_i(j) = \vec{x}(i+j-1)$  and  $\vec{z}_i(j+1) = \vec{x}(i+j)$ . Since  $i+j \leq i+k-i = k$ , we have  $\vec{x}(i+j-1) > \vec{x}(i+j)$  and hence  $\vec{z}_i(j) > \vec{z}_i(j+1)$ . Thus,  $k(\vec{z}_i) \geq k-i+1$ . Note  $\vec{z}_i(k-i+1) = \vec{x}(i+(k-i+1)-1) = \vec{x}(k)$  and  $\vec{z}_i(k-i+2) = \vec{x}(i+(k-i+2)-1) = \vec{x}(k+1)$ . By definition,  $\vec{x}(k) \leq \vec{x}(k+1)$  and hence  $\vec{z}_i(k-i+1) \leq \vec{z}_i(k-i+2)$ . Therefore,  $k(\vec{z}_i) = k-i+1$ . As we have already seen,  $\vec{z}_i(k(\vec{z}_i)+1) = \vec{z}_i(k-i+2) = \vec{x}(k+1)$ . ⊣ (Claim 5)

**Claim 6.** For every  $i \in \{1, \dots, l-1\}$ , if  $\vec{x}(i) > \vec{x}(i+1)+1$ , then  $l(\vec{z}_i) = l-i+1$  and hence  $\vec{z}_i(l(\vec{z}_i)+2) = \vec{x}(l+2)$ .

⊢ Since  $\vec{x}(i) > \vec{x}(i+1) + 1$ , we have  $\vec{z}_i(1) > \vec{z}_i(2)$ . Since  $i < l$  and  $\vec{x}(i) > \vec{x}(i+1) + 1$ , we have  $\vec{z}_i(2) = \vec{x}(i+1) > \vec{x}(i+2) + 1 = \vec{z}_i(3) + 1$ . Thus,  $l(\vec{z}_i) \geq 2$ .

Let  $j \in \{2, \dots, l-i\}$ . Then,  $\vec{z}_i(j) = \vec{x}(i+j-1)$ ,  $\vec{z}_i(j+1) = \vec{x}(i+j)$ , and  $\vec{z}_i(j+2) = \vec{x}(i+j+1)$ . Note  $i+j-1 \leq i+(l-i)-1 = l-1$  and hence  $i+j \leq l < k$ . So, either  $\vec{x}(i+j-1) = \vec{x}(i+j)+1$  or  $\vec{x}(i+j) > \vec{x}(i+j+1)+1$ . Thus, either  $\vec{z}_i(j) = \vec{z}_i(j+1)+1$  or  $\vec{z}_i(j+1) > \vec{z}_i(j+2)+1$ . Hence  $l(\vec{z}_i) \geq l-i+1$ . If  $l = k-1$ , then by Claim 5,  $k(\vec{z}_i) = k-i+1$  and hence  $l(\vec{z}_i) \leq k(\vec{z}_i) - 1 = k-i = l-i+1$ . Thus,  $l(\vec{z}_i) = l-i+1$ . If  $l < k-1$ , then we have both  $\vec{x}(l) > \vec{x}(l+1) + 1$  and  $\vec{x}(l+1) = \vec{x}(l+2) + 1$ . Note  $l-i+1 \geq l-(l-1)+1 = 2$ . So,  $\vec{z}_i(l-i+1) = \vec{x}(l)$ ,  $\vec{z}_i(l-i+2) = \vec{x}(l+1)$  and  $\vec{z}_i(l-i+3) = \vec{x}(l+2)$ . Thus,  $\vec{z}_i(l-i+1) > \vec{z}_i(l-i+2)+1$  and  $\vec{z}_i(l-i+2) = \vec{z}_i(l-i+3)+1$ . So,  $l(\vec{z}_i) = l-i+1$ . ⊣ (Claim 6)



**Claim 7.** For every  $i \in \{1, \dots, l-1\}$ , if  $\vec{x}(i) > \vec{x}(i+1)+1$ , then  $f(\vec{z}_i) = f(\vec{x})$ .

$\vdash$  By Claim 5,  $\vec{z}_i(k(\vec{z}_i) + 1) = \vec{x}(k+1)$ . By Claim 6,  $\vec{z}_i(l(\vec{z}_i) + 2) = \vec{x}(l+2)$ . Therefore,

$$\begin{aligned} f(\vec{z}_i) &= \max\{\vec{z}_i(l(\vec{z}_i) + 2), \vec{z}_i(k(\vec{z}_i) + 1)\} \\ &= \max\{\vec{x}(l+2), \vec{x}(k+1)\} = f(\vec{x}) \end{aligned}$$

$\dashv$  (Claim 7)

**Case 1.**  $m+2 \leq l$ .

Then, by Claim 7,  $\vec{y}(m) = \vec{y}(m+1) = f(\vec{x})$ . By Claim 4,  $k(\vec{y}) \geq m-1$  and  $l(\vec{y}) \geq m-2$ . By Lemma 3.7,  $f(\vec{y}) = f(\vec{x})$ .

**Case 2.**  $m+1 = l$ .

Then by Claim 7,  $\vec{y}(m) = f(\vec{x})$ .

**Subcase 2.1.**  $l+1 = k$ .

Then, we have  $f(\vec{x}) = \vec{x}(k+1)$  and

$$\vec{y}(m+2) = \vec{y}(l+1) = \vec{y}(k) = \vec{x}(k+1) = f(\vec{x})$$

Consider  $\vec{z}_l$ . We have

$$\begin{aligned} \vec{z}_l(1) &= \vec{x}(l) - 1 \\ \vec{z}_l(2) &= \vec{x}(l+1) = \vec{x}(k) \\ \vec{z}_l(3) &= \vec{x}(l+2) = \vec{x}(k+1) \end{aligned}$$

By the definition of  $k$ ,  $\vec{x}(k) \leq \vec{x}(k+1)$ . Thus,  $\vec{z}_l(2) \leq \vec{z}_l(3)$ . By Lemma 3.6,  $f(\vec{z}_l)$  is either  $\vec{z}_l(2)$  or  $\vec{z}_l(3)$ , i.e. either  $\vec{x}(k)$  or  $\vec{x}(k+1)$ . In particular, we have  $f(\vec{z}_l) \leq \vec{x}(k+1)$  and hence  $\vec{y}(m+1) = \vec{y}(m) \leq \vec{x}(k+1)$ . Therefore, we have  $k(\vec{y}) \geq m-1$ ,  $l(\vec{y}) \geq m-2$ ,  $\vec{y}(m) = \vec{y}(m+2) = \vec{x}(k+1)$ , and  $\vec{y}(m+1) \leq \vec{y}(m)$ . By Lemma 3.8, we have  $f(\vec{y}) = \vec{x}(k+1) = f(\vec{x})$ .

**Subcase 2.2.**  $l < k-1$ .

Then, we have  $\vec{x}(l) > \vec{x}(l+1) + 1$ . By Claim 5,  $k(\vec{z}_l) = k-l+1$  and  $\vec{z}_l(k(\vec{z}_l) + 1) = \vec{x}(k+1)$ .

**Subsubcase 2.2.1.**  $\vec{x}(l+2) \leq \vec{x}(k+1)$

Then,  $\vec{z}_l(3) = \vec{x}(l+3-1) = \vec{x}(l+2) \leq \vec{x}(k+1)$  By Lemma 3.3,  $f(\vec{z}_l) = \vec{z}_l(k(\vec{z}_l) + 1) = \vec{x}(k+1)$ . Therefore, we have  $\vec{y}(m) = \vec{y}(m+1) = \vec{x}(k+1)$ . Since  $k(\vec{y}) \geq m-1$  and  $l(\vec{y}) \geq m-2$ , by Lemma 3.7, we have  $f(\vec{y}) = \vec{y}(m) = \vec{x}(k+1) = f(\vec{x})$ .

**Subsubcase 2.2.2.**  $\vec{x}(l+2) > \vec{x}(k+2)$

Recall  $k(\vec{z}_l) = k - l + 1$ . Note  $k - l + 1 \geq 1 + 1 = 2$ . By Lemma 3.9,

$$\begin{aligned} f(\vec{z}_l) &\leq \max\{\vec{z}_l(3), \vec{z}_l(k(\vec{z}) + 1)\} \\ &= \max\{\vec{x}(l + 2), \vec{x}(k + 2)\} = \vec{x}(l + 2) \end{aligned}$$

Therefore,  $\vec{y}(m + 1) = \vec{y}(l) \leq \vec{x}(l + 2)$ . Since  $l < k - 1$ , we have  $\vec{x}(l + 1) = \vec{x}(l + 2) + 1$ . So,  $\vec{z}_{l+1}(1) = \vec{x}(l + 1) - 1 = \vec{x}(l + 2) = \vec{z}_{l+1}(2)$ . Therefore,  $f(\vec{z}_{l+1}) = \vec{z}_{l+1}(2) = \vec{x}(l + 2)$ . Hence,

$$\begin{aligned} \vec{y}(m) &= \vec{x}(l + 2) \\ \vec{y}(m + 1) &\leq \vec{x}(l + 2) \\ \vec{y}(m + 2) &= \vec{x}(l + 2) \end{aligned}$$

By Lemma 3.8, we have  $f(\vec{y}) = \vec{y}(m) = \vec{x}(l + 2) = f(\vec{x})$ .

**Case 3.**  $m = l$

**Subcase 3.1.**  $l = k - 1$ .

**Claim 8.**  $\vec{y}(l) = f(\vec{z}_l)$  is either  $\vec{x}(k)$  or  $\vec{x}(k + 1)$ . In particular,  $\vec{y}(l) = f(\vec{z}_l) \leq \vec{x}(k + 1) = f(\vec{x})$ .

└ Since we assumed  $l = k - 1$ ,  $\vec{y}(l + 1) = \vec{y}(k) = \vec{x}(k + 1)$ . Note

$$\begin{aligned} \vec{z}_l(2) &= \vec{x}(l + 1) = \vec{x}(k) \\ &\leq \vec{x}(k + 1) = \vec{x}(l + 2) = \vec{z}_l(3) \end{aligned}$$

Thus, by Lemma 3.6,  $f(\vec{z}_l)$  is either  $\vec{x}(k)$  or  $\vec{x}(k + 1)$ . ┘ (Claim 8)

**Claim 9.**  $k(\vec{y}) \leq l$ .

└ Because

$$\vec{y}(l) \leq \vec{x}(k + 1) = \vec{y}(k) = \vec{y}(l + 1)$$

┘ (Claim 9)

**Subsubcase 3.1.1.**  $l = 1$

Then, by Claim 8,  $\vec{y}(1) = \vec{y}(l) \leq \vec{x}(k + 1)$ . By Claim 1,  $\vec{y}(2) = \vec{x}(k) = \vec{x}(k + 1)$ . So,  $f(\vec{y}) = \vec{x}(k + 1) = f(\vec{x})$ .

**Subsubcase 3.1.2.**  $l \geq 2$  and  $\vec{x}(l) \leq \vec{x}(k + 1)$ .

Recall that by Claim 8,  $\vec{y}(l)$  is either  $\vec{x}(k + 1)$  or  $\vec{x}(k)$ . If  $\vec{y}(l) = \vec{x}(k + 1)$ , then since  $\vec{y}(l - 1) = \vec{x}(l) \leq \vec{x}(k + 1) = \vec{y}(l)$ , we have  $k(\vec{y}) \leq l - 1$ . Since we also know  $k(\vec{y}) \geq m - 1 = l - 1$  by Claim 4, we have  $k(\vec{y}) = l - 1$ . Since  $l - 2 = m - 2 \leq l(\vec{y}) \leq k(\vec{y}) - 1 = l - 2$ , we have  $l(\vec{y}) = l - 2$ . Therefore,  $f(\vec{y}) = \vec{x}(k + 1) = f(\vec{x})$ .

Suppose  $\vec{y}(l) = \vec{x}(k)$ . Since  $l < k$ , we have  $\vec{y}(l - 1) = \vec{x}(l) > \vec{x}(k) = \vec{y}(l)$ . Thus,  $k(\vec{y}) \geq l$ . By Claim 9,  $k(\vec{y}) \leq l$  and hence  $k(\vec{y}) = l$ . Note that  $l(\vec{y}) \geq m - 2 = l - 2 = k(\vec{y}) - 2$ . By Lemma 3.5,  $f(\vec{y}) = \vec{y}(k(\vec{y}) + 1) = \vec{y}(l + 1) = \vec{x}(k + 1) = f(\vec{x})$ . Therefore, in either case, we get  $f(\vec{y}) = f(\vec{x})$ .

**Subsubcase 3.1.3.**  $l \geq 2$  and  $\vec{x}(l) > \vec{x}(k+1)$ .

We have  $\vec{y}(l-1) = \vec{x}(l) > \vec{x}(k+1) = \vec{y}(l)$ . Hence,  $k(\vec{y}) \geq l$ . By Claim 9, we have  $k(\vec{y}) = l$ . Note that  $l(\vec{y}) \geq m-2 = l-2 = k(\vec{y})-2$ . By Lemma 3.5, we have  $f(\vec{y}) = \vec{y}(k(\vec{y})+1) = \vec{y}(l+1) = \vec{x}(k+1) = f(\vec{x})$ .

**Subcase 3.2.**  $l < k-1$ .

**Claim 10.**  $\vec{y}(l+1) = \vec{x}(l+2)$

└ Since  $l < k-1$ , we have  $\vec{x}(l) > \vec{x}(l+1)+1$  and  $\vec{x}(l+1) = \vec{x}(l+2)+1$ .  
Thus,  $\vec{y}(l+1) = \vec{x}(l+2)$ . ┘ (Claim 10)

**Subsubcase 3.2.1.**  $\vec{x}(l+2) \geq \vec{x}(k+1)$

By Claim 5,  $k(\vec{z}_l) = k-l+1$  and  $\vec{z}_l(k(\vec{z}_l)+1) = \vec{x}(k+1)$ . We have

$$\begin{aligned} l &< k-1 \\ 1 &< k-l \\ 2 &< k-l+1 \end{aligned}$$

By Lemma 3.9,

$$\begin{aligned} \vec{y}(l) &= f(\vec{z}_l) \leq \max\{\vec{z}_l(3), \vec{z}_l(k(\vec{z}_l)+1)\} \\ &= \max\{\vec{x}(l+2), \vec{x}(k+1)\} = \vec{x}(l+2) = f(\vec{x}) \end{aligned}$$

Then,  $\vec{y}(l) \leq \vec{x}(l+2) = \vec{y}(l+1)$ . So, we have  $k(\vec{y}) \leq l$ . If  $l = 1$ , then clearly  $k(\vec{y}) = 1 = l$ . If  $l \geq 2$ , then since  $l < l+2 \leq k$ ,

$$\vec{y}(l-1) = \vec{x}(l) > \vec{x}(l+2) \geq \vec{y}(l)$$

So,  $k(\vec{y}) \geq l$  and hence  $k(\vec{y}) = l$ . Therefore, in either case, we get  $k(\vec{y}) = l$ .

We also have  $l(\vec{y}) \geq m-2 = l-2 = k(\vec{y})-2$ . By Lemma 3.5,  $f(\vec{y}) = \vec{y}(k(\vec{y})+1) = \vec{y}(l+1) = \vec{x}(l+2) = f(\vec{x})$ .

Now we concentrate on the case  $\vec{x}(l+2) < \vec{x}(k+1)$ .

**Claim 11.** If  $\vec{x}(l+2) < \vec{x}(k+1)$ , then  $\vec{y}(l) = \vec{x}(k+1) = f(\vec{x})$ .

└ Since  $l < k-1$ , we have  $\vec{x}(l) > \vec{x}(l+1)+1$ . By Claim 5,  $\vec{z}_l(k(\vec{z}_l)+1) = \vec{x}(k+1)$ . In addition,  $\vec{z}_l(3) = \vec{x}(l+2) < \vec{x}(k+1)$ . By Lemma 3.3,  $\vec{y}(l) = f(\vec{z}_l) = \vec{x}(k+1)$ . ┘ (Claim 11)

**Subsubcase 3.2.2.**  $\vec{x}(l+2) < \vec{x}(k+1)$  and  $k(\vec{y}) = l-1$ .

Note  $l(\vec{y}) \geq l-2 \geq k(\vec{y})-1$ . By Lemma 3.2,  $f(\vec{y}) = \vec{y}(k(\vec{y})+1) = \vec{y}(l) = f(\vec{x})$ .

**Subsubcase 3.2.3.**  $\vec{x}(l+2) < \vec{x}(k+1)$  and  $k(\vec{y}) \geq l$ .

Let  $p$  be the least such that  $l+1 \leq p < k$  and  $\vec{x}(p) > \vec{x}(p+1)+1$  if exists. Otherwise, let  $p = k$ .

**Claim 12.**  $p \geq l + 2$ .

$\vdash$  By definition,  $p \geq l + 1$ . Since  $l < k - 1$ ,  $\vec{x}(l + 1) = \vec{x}(l + 2) + 1$ . So,  
 $p \geq l + 2$ .  $\dashv$  (Claim 12)

**Claim 13.**  $k(\vec{y}) = p - 1$ .

$\vdash$  By assumption, we have  $k(\vec{y}) \geq l$ .

**Subclaim 13.1.**  $\vec{y}(l) > \vec{y}(l + 1)$ . In particular,  $k(\vec{y}) \geq l + 1$ .

$\vdash$  By Claim 11,  $\vec{y}(l) = \vec{x}(k + 1)$ . By assumption,  $\vec{x}(k + 1) > \vec{x}(l + 2) = \vec{y}(l + 1)$ .  $\dashv$  (Subclaim 13.1)

**Subclaim 13.2.** For every  $i \in \{l + 1, \dots, p - 2\}$ ,  $\vec{y}(i) = \vec{y}(i + 1) + 1$ . In particular,  $k(\vec{y}) \geq p - 1$ .

$\vdash$  By the definition of  $p$ , since  $l < i < i + 1 < p$ ,  $\vec{x}(i) = \vec{x}(i + 1) + 1$  and  $\vec{x}(i + 1) = \vec{x}(i + 2) + 1$ . Thus,  $\vec{y}(i) = \vec{x}(i + 1)$  and  $\vec{y}(i + 1) = \vec{x}(i + 2)$ . So,  
 $\vec{y}(i) = \vec{x}(i + 1) = \vec{x}(i + 2) + 1 = \vec{y}(i + 1) + 1$ .  $\dashv$  (Subclaim 13.2)

**Subclaim 13.3.**  $\vec{y}(p) = \vec{x}(k + 1)$ .

$\vdash$  If  $p = k$ , then we have  $\vec{y}(p) = \vec{y}(k) = \vec{x}(k + 1)$ . If  $p = k - 1$ , then we have  $\vec{z}_p(1) = \vec{x}(p) - 1 > \vec{x}(p + 1) = \vec{z}_p(2)$  and  $\vec{z}_p(2) = \vec{x}(p + 1) = \vec{x}(k) \leq \vec{x}(k + 1) = \vec{z}_p(3)$ . So,  $\vec{y}(p) = f(\vec{z}_p) = \vec{x}(k + 1)$ .

Suppose  $p \leq k - 2$ . By Claim 5,  $\vec{z}_p(k(\vec{z}_p) + 1) = \vec{x}(k + 1)$ . Note  $l + 2 \leq p \leq p + 2 \leq k$ , so,  $\vec{z}_p(3) = \vec{x}(p + 2) \leq \vec{x}(l + 2) < \vec{x}(k + 1) = \vec{z}_p(k(\vec{z}_p) + 1)$ . By Lemma 3.3,  $y(p) = f(\vec{z}_p) = \vec{z}_p(k(\vec{z}_p) + 1) = \vec{x}(k + 1)$ .  $\dashv$  (Subclaim 13.3)

**Subclaim 13.4.**  $\vec{y}(p - 1) < \vec{x}(k + 1) = \vec{y}(p)$ . In particular,  $k(\vec{z}) \leq p - 1$ .

$\vdash$  Since  $p \geq l + 2$ , we have  $p - 1 \geq l + 1 > l$ . By the definition of  $p$ , we have  $\vec{x}(p - 1) = \vec{x}(p) + 1$  and hence  $\vec{y}(p - 1) = \vec{x}(p)$ . Since  $l + 2 \leq p \leq k$ , we have  $\vec{x}(p) \leq \vec{x}(l + 2)$ . By assumption,  $\vec{x}(l + 2) < \vec{x}(k + 1)$ . Therefore,  
 $\vec{y}(p - 1) < \vec{x}(k + 1)$ .  $\dashv$  (Subclaim 13.4)

By Subclaim 13.2 and Subclaim 13.4, we have  $k(\vec{y}) = p - 1$ .  $\dashv$  (Claim 13)

If  $l(\vec{y}) = k(\vec{y}) - 1$ , then clearly  $f(\vec{y}) = \vec{y}(k(\vec{y}) + 1) = \vec{y}(p) = \vec{x}(k + 1) = f(\vec{x})$ . Suppose  $l(\vec{y}) < k(\vec{y}) - 1$ . By Claim 4,  $l(\vec{y}) \geq m - 1 = l - 1$ . Thus,  $l + 1 \leq l(\vec{y}) + 2$ . Recall  $\vec{y}(l + 1) = \vec{x}(l + 2) < \vec{x}(k + 1) = \vec{y}(p) = \vec{y}(k(\vec{y}) + 1)$ . By Lemma 3.2,  $f(\vec{y}) = \vec{y}(k(\vec{y}) + 1) = f(\vec{x})$ .  $\square$ (Lemma 3.10)

**Lemma 3.11.** For every  $\vec{x} \in \mathbb{Z}^n$ ,  $f(\vec{x})$  satisfies the tarai recurrence.

*Proof.* By Lemma 3.10, we may assume that  $k(\vec{x}) = n$ , i.e.  $\vec{x}(1) > \vec{x}(2) > \dots > \vec{x}(n)$ . For every  $i = 1, \dots, n$ , define  $\vec{z}_i = \sigma(r^{i-1}(\vec{x}))$  and  $\vec{y}(i) = f(\vec{z}_i)$ . We need to show that  $f(\vec{y}) = f(\vec{x}) = \vec{x}(1)$ .

**Claim 1.**  $\vec{y}(1) < \vec{x}(1)$ .

⊢ We have  $\vec{z}_1(1) = \vec{x}(1) - 1$  and  $\vec{z}_1(i) = \vec{x}(i)$  for every  $i = 2, \dots, n$ . Then, clearly we have  $\vec{y}(1) = f(\vec{z}_1) \leq \max \vec{z}_1 = \vec{x}(1) - 1$ . ⊣ (Claim 1)

**Claim 2.**  $\vec{y}(n) = \vec{x}(1)$

⊢ Since  $\vec{z}_n(1) = \vec{x}(n)$  and  $\vec{z}_n(2) = \vec{x}(1)$ , we have  $\vec{y}(n) = f(\vec{z}_n) = \vec{z}(2) = \vec{x}(1)$ . ⊣ (Claim 2)

**Claim 3.** For every  $i = 2, \dots, n - 1$ , either  $\vec{y}(i) = \vec{x}(1)$  or  $\vec{y}(i) = \vec{x}(i + 1)$ .

⊢ Note

$$\begin{aligned}\vec{z}_i(1) &= \vec{x}(i) - 1 \\ \vec{z}_i(j) &= \vec{x}(j + i - 1) \text{ (for all } j = 2, \dots, n - i + 1) \\ \vec{z}_i(n - i + 2) &= \vec{x}(1)\end{aligned}$$

If  $\vec{x}(i) - 1 = \vec{x}(i + 1)$ , then we have  $\vec{z}_i(1) = \vec{z}_i(2)$  and hence  $\vec{y}(i) = f(\vec{z}_i) = \vec{z}_i(2) = \vec{x}(i + 1)$ .

Suppose  $\vec{x}(i) - 1 > \vec{x}(i + 1)$ . Then, for every  $j = 2, \dots, n - i$ , we have  $\vec{z}_i(j) = \vec{x}(j + i - 1) > \vec{x}(j + i) = \vec{z}_i(j + 1)$ . Moreover,  $\vec{z}_i(n - i + 1) = \vec{x}(n) < \vec{x}(1) = \vec{z}_i(n - i + 2)$ . So,  $k(\vec{z}_i) = n - i + 1$  and  $\vec{z}_i(k(\vec{z}_i) + 1) = \vec{z}_i(n - i + 2) = \vec{x}(1)$ . Then, since  $\vec{z}_i(2) = \vec{x}(i + 1) < \vec{x}(1) = \vec{z}_i(k(\vec{z}_i) + 1)$ , by Lemma 3.4, we have  $\vec{y}(i) = f(\vec{z}_i) = \vec{z}_i(k(\vec{z}_i) + 1) = \vec{x}(1)$ . ⊣ (Claim 3)

Let  $m$  be the least such that  $\vec{y}(m + 1) = \vec{x}(1)$ . Since  $\vec{y}(n) = \vec{x}(1)$ , there is such an  $m \leq n - 1$ . If  $m = 1$ , then we have  $\vec{y}(1) < \vec{x}(1) = \vec{y}(2)$  and hence  $f(\vec{y}) = \vec{y}(2) = \vec{x}(1)$ .

Suppose  $m > 1$ . Then for every  $i = 2, \dots, m$ , by Claim 3, we have  $\vec{y}(i) = \vec{x}(i + 1)$ . Hence, for every  $i = 2, \dots, m - 1$ , we have  $\vec{y}(i) > \vec{y}(i + 1)$ . We also have  $\vec{y}(m) = \vec{x}(m + 1) < \vec{x}(1) = \vec{y}(m + 1)$ . Therefore,  $k(\vec{y}) = m$  and  $\vec{y}(m + 1) = \vec{x}(1)$ . In particular,  $\vec{y}(2) = \vec{x}(3) < \vec{x}(1) = \vec{y}(m + 1) = \vec{y}(k(\vec{y}) + 1)$ . By Lemma 3.4, we have  $f(\vec{y}) = \vec{y}(k(\vec{y}) + 1) = \vec{x}(1)$ . □(Lemma 3.11)

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